Finite Models

Learning objectives

- Understand goals and implications of finite state abstraction
- Learn how to model program control flow with graphs
- Learn how to model the software system structure with **call graphs**
- Learn how to model finite state behavior with **finite state machines**

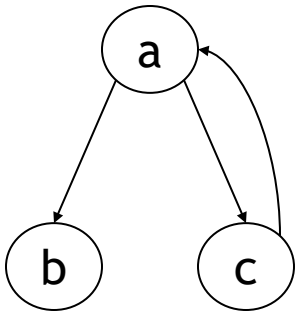
Properties of Models

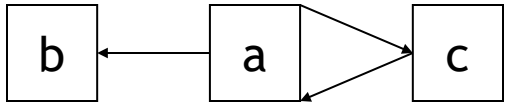
- **Compact:** representable and manipulable in a reasonably compact form
  - What is *reasonably compact* depends largely on how the model will be used
- **Predictive:** must represent some salient characteristics of the modeled artifact well enough to distinguish between *good* and *bad* outcomes of analysis
  - no single model represents all characteristics well enough to be useful for all kinds of analysis
- **Semantically meaningful:** it is usually necessary to interpret analysis results in a way that permits diagnosis of the causes of failure
- **Sufficiently general:** models intended for analysis of some important characteristic must be general enough for practical use in the intended domain of application

Graph Representations: directed graphs

- Directed graph:
  - N (set of nodes)
  - E (relation on the set of nodes) edges

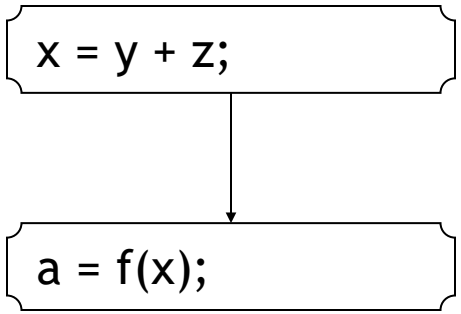
Nodes: {a, b, c}
Edges: {(a,b), (a, c), (c, a)}





Graph Representations: labels and code

- We can label nodes with the names or descriptions of the entities they represent.
  - If nodes a and b represent program regions containing assignment statements, we might draw the two nodes and an edge (a,b) connecting them in this way:

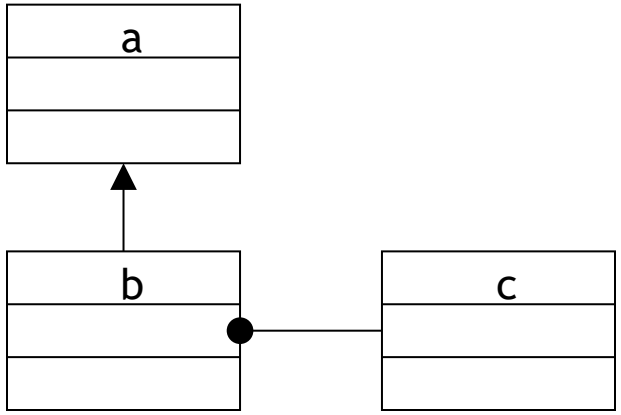


Multidimensional Graph Representations

- Sometimes we draw a single diagram to represent more than one directed graph, drawing the shared nodes only once
  - class B extends (is a subclass of) class A
  - class B has a field that is an object of type C

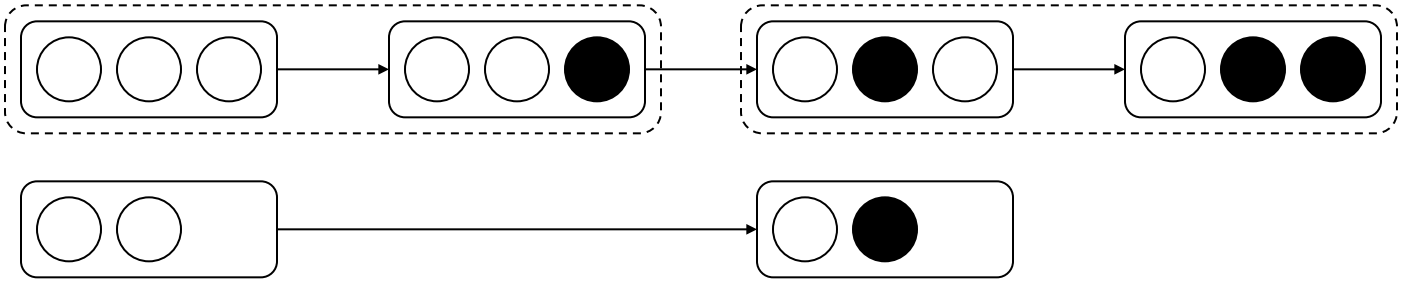
extends relation
NODES = {A, B, C}
EDGES = {(A,B)}

includes relation
NODES = {A, B, C}
EDGES = {(B,C)}

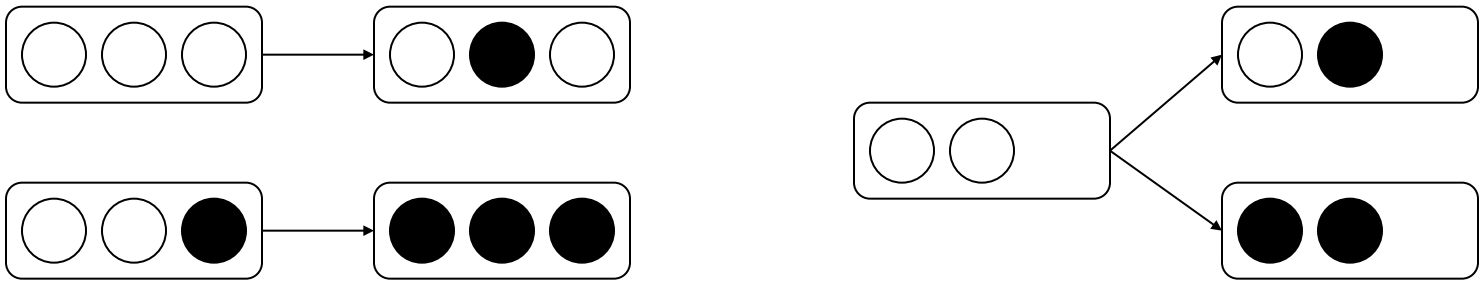


Finite Abstraction of Behavior

An abstraction function suppresses some details of program execution. it lumps together execution states that differ with respect to the suppressed details but are otherwise identical.



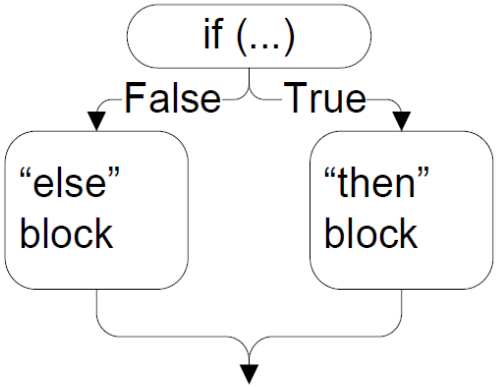
Program execution states with different successor states can been merge:

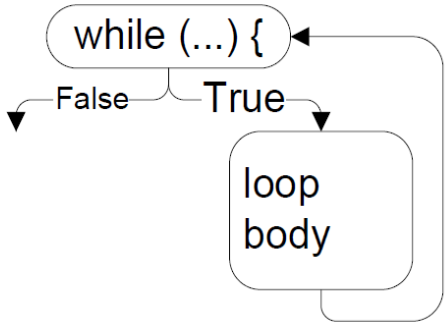


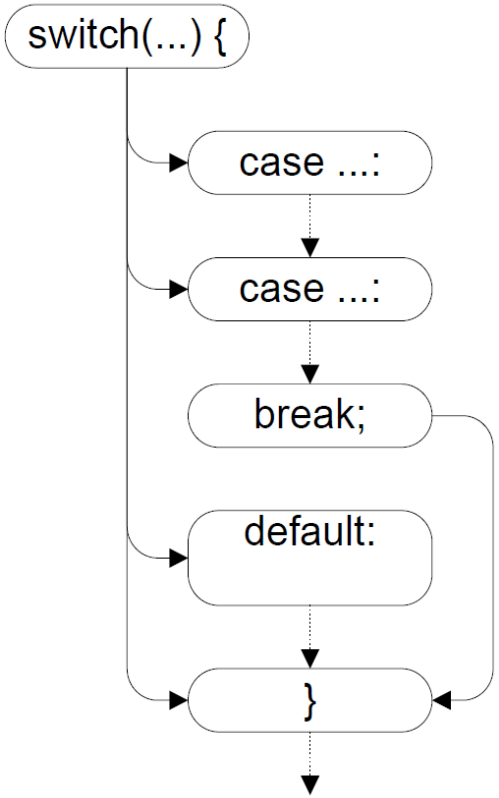
(Intraprocedural) Control Flow Graph (CFG)

- nodes = regions of source code (basic blocks)
  - Basic block = maximal program region with a single entry and single exit point
  - Often statements are grouped in single regions to get a compact model
  - Sometime single statements are broken into more than one node to model control flow within the statement
- directed edges = possibility that program execution proceeds from the end of one region directly to the beginning of another

Building blocks for CFG





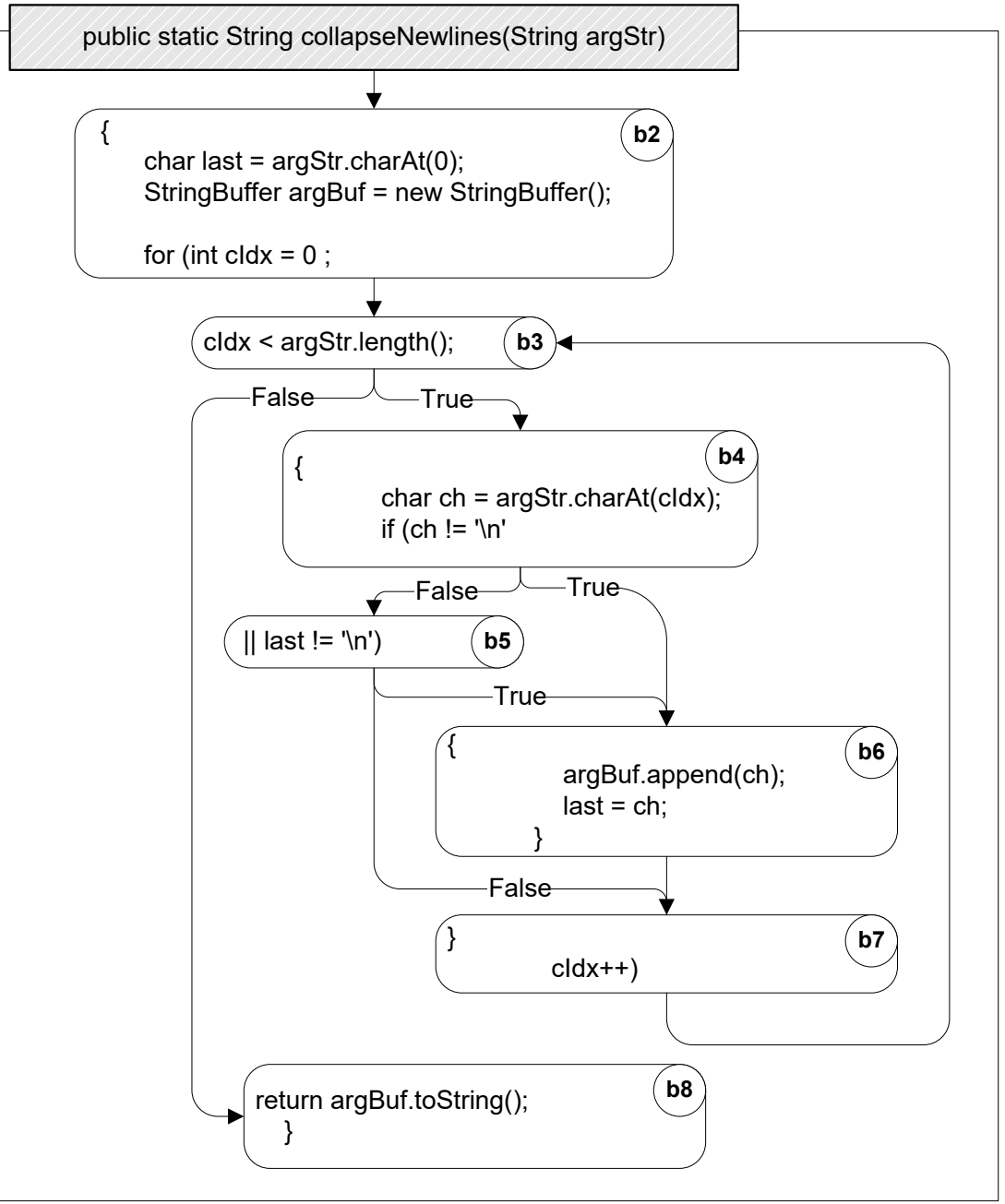


Example of Control Flow Graph

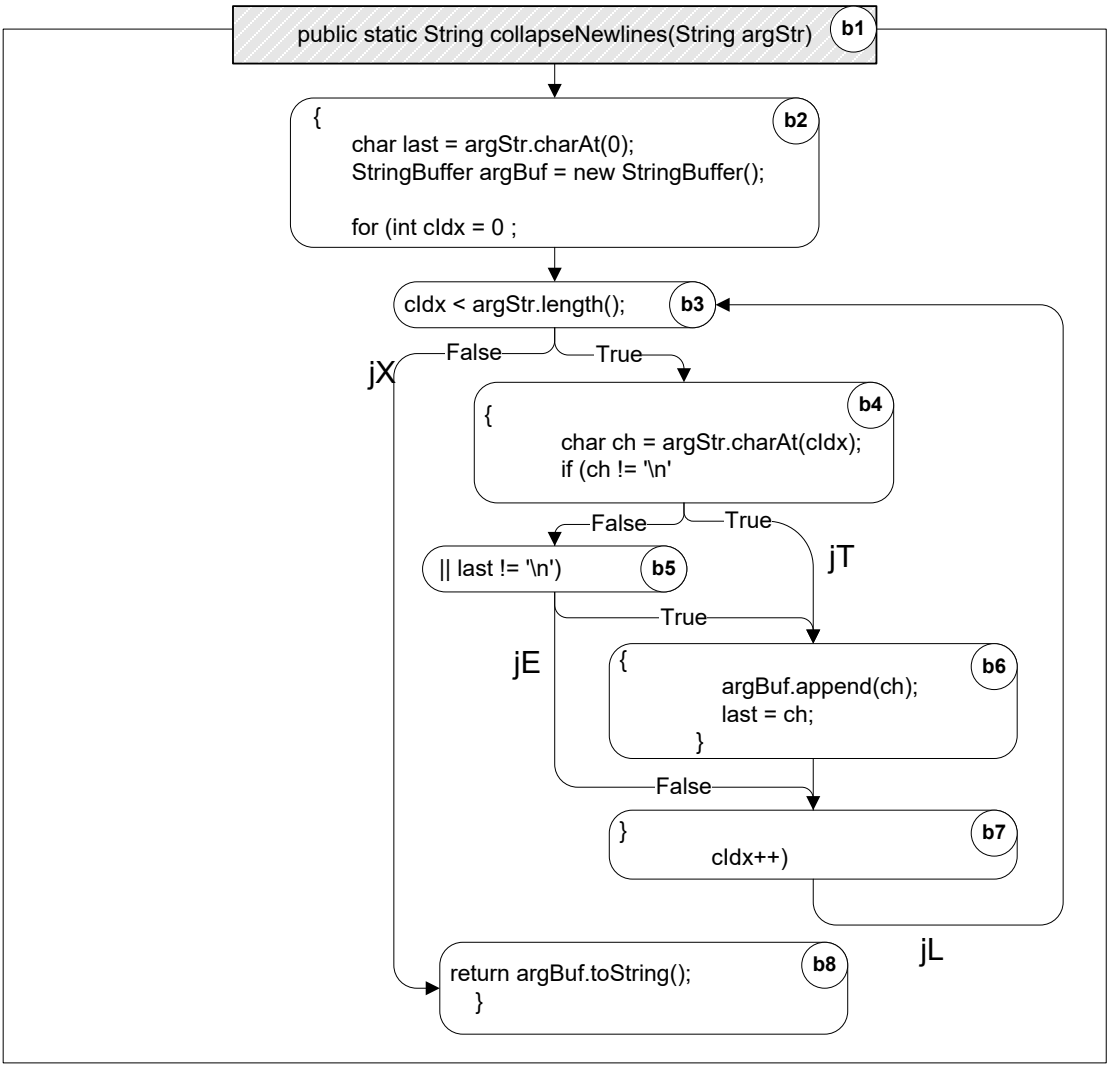
```
public static String collapseNewlines(String argStr)
   {
      char last = argStr.charAt(0);
      StringBuffer argBuf = new StringBuffer();

      for (int cIdx = 0 ; cIdx < argStr.length(); cIdx++)
      {
         char ch = argStr.charAt(cIdx);
         if (ch != '\n' || last != '\n')
         {
            argBuf.append(ch);
            last = ch;
         }
      }

      return argBuf.toString();
   }
```



Linear Code Sequence and Jump (LCSJ)



Essentially subpaths of the control flow graph from one branch to another

From	Sequence of basic blocs	To
Entry	b1 b2 b3	jX
Entry	b1 b2 b3 b4	jT
Entry	b1 b2 b3 b4 b5	jE
Entry	b1 b2 b3 b4 b5 b6 b7	jL
jX	b8	ret
jL	b3 b4	jT
jL	b3 b4 b5	jE
jL	b3 b4 b5 b6 b7	jL

Pen and paper exercise

Draw the control flow graph for the following code.

```
int evensum(int i)
{
  int sum = 0;

  while (i <= 10) {
    if (i/2 == 0)
      sum = sum + i;
    i++;
  }
  return sum;

}
```

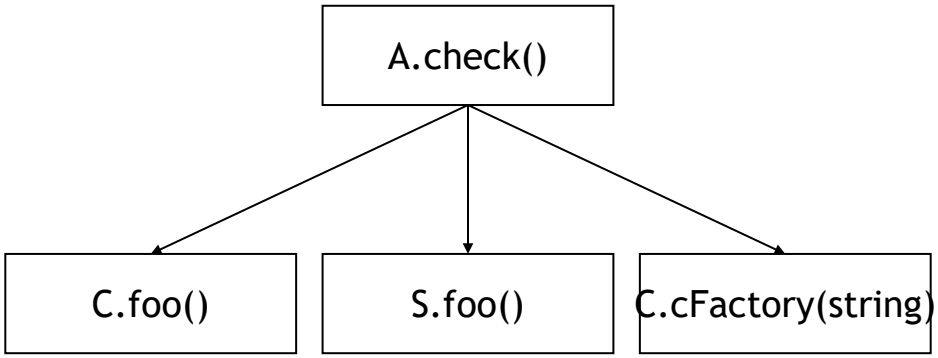
Interprocedural control flow graph

- Call graphs
  - Nodes represent procedures
    - Methods
    - C functions
    - ...
  - Edges represent *calls* relation

Overestimating the *calls* relation

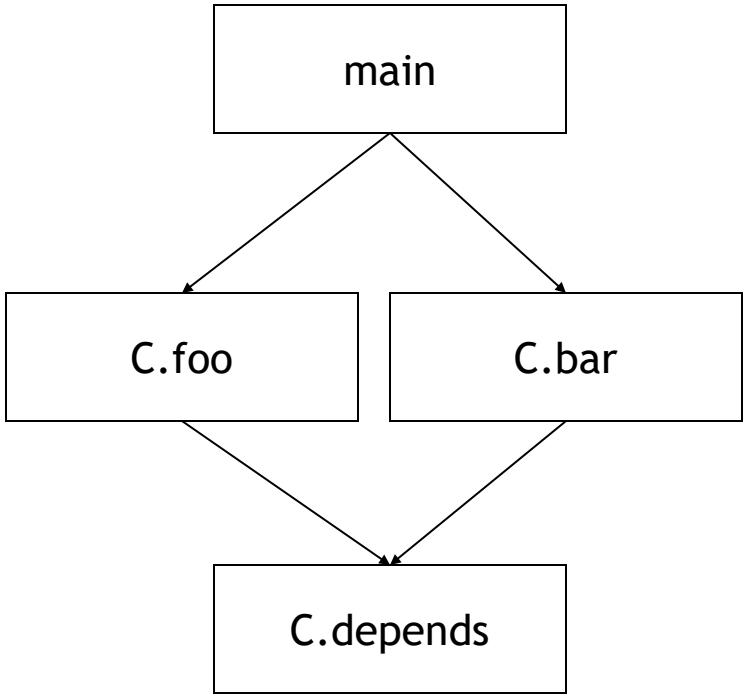
```
public class C {
   public static C cFactory(String kind) {
      if (kind == "C") return new C();
      if (kind == "S") return new S();
      return null;
   }
   void foo() {
      System.out.println("You called the parent's method");
   }
   public static void main(String args[]) {
      (new A()).check();
   }
}
class S extends C {
   void foo() {
      System.out.println("You called the child's method");
   }
}
class A {
   void check() {
      C myC = C.cFactory("S");
      myC.foo();
   }
}
```

The static call graph includes calls through dynamic bindings that *never* occur in execution.



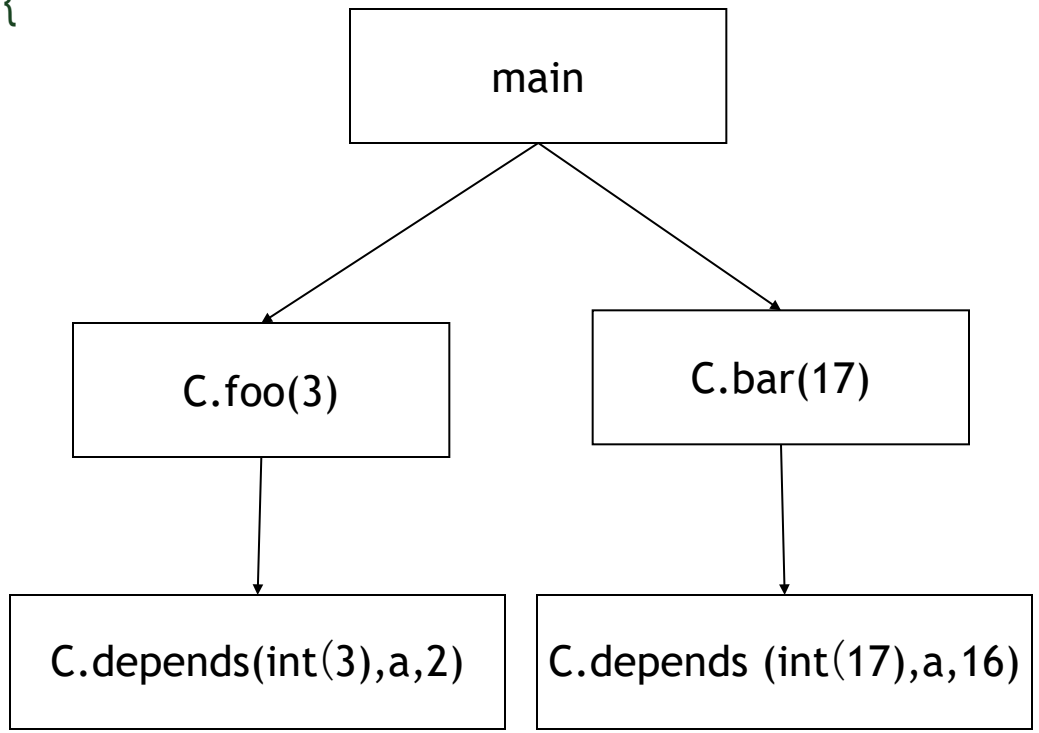
Contex Insensitive Call graphs

```
public class Context {
    public static void main(String args[]) {
     Context c = new Context();
     c.foo(3);
     c.bar(17);
    }
    void foo(int n) {
     int[]  myArray = new int[ n ];
     depends( myArray, 2) ;
    }
    void bar(int n) {
     int[]  myArray = new int[ n ];
     depends( myArray, 16) ;
    }
    void depends( int[] a, int n ) {
     a[n] = 42;
    }
}
```

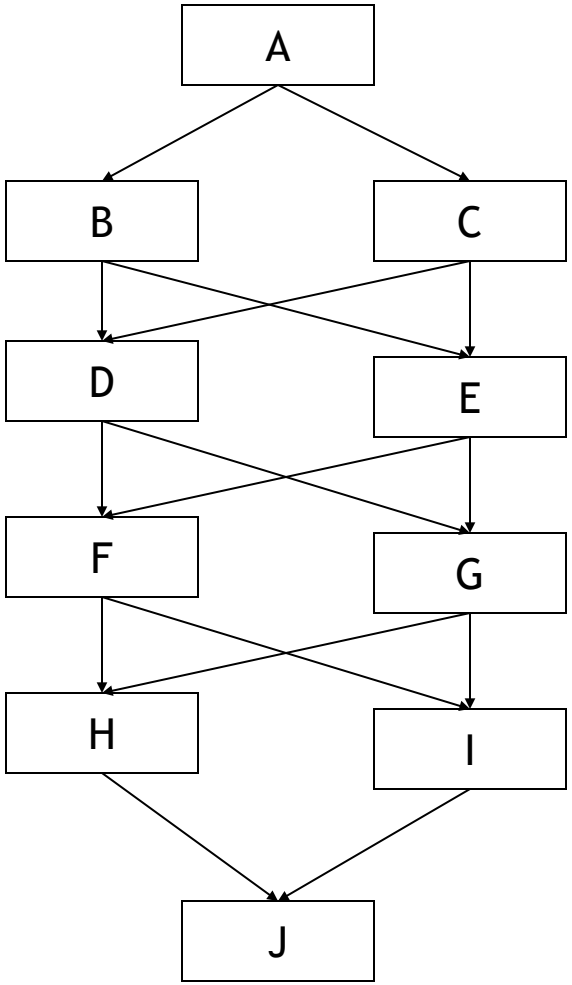


Contex Sensitive Call graphs

```
public class Context {
    public static void main(String args[]) {
     Context c = new Context();
     c.foo(3);
     c.bar(17);
    }
    void foo(int n) {
     int[]  myArray = new int[ n ];
     depends( myArray, 2) ;
    }
    void bar(int n) {
     int[]  myArray = new int[ n ];
     depends( myArray, 16) ;
    }
    void depends( int[] a, int n ) {
     a[n] = 42;
    }
}
```



Context Sensitive CFG exponential growth



1 context A

2 contexts AB AC

4 contexts ABD ABE ACD ACE

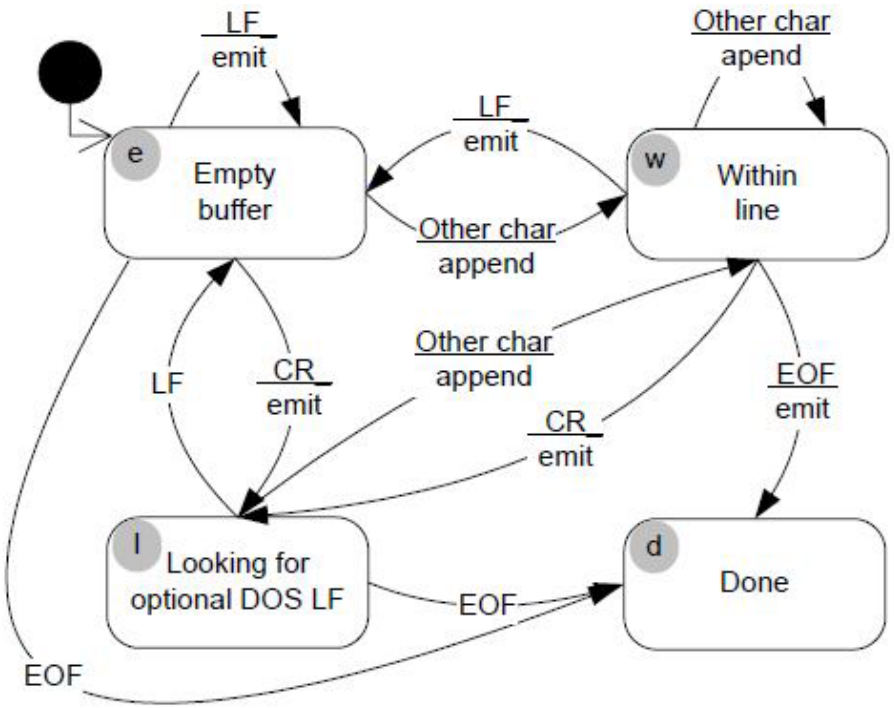
8 contexts …

16 calling contexts …

Finite state machines

- finite set of states (nodes) with transitions among states (edges)
- $\frac{\text{Event}}{\text{Respond}}$ means "event, respond" (e.g., $\frac{\text{LF}}{\text{emit}}$ means "receiving *LF*, responding *emit*".)

Example of graph representation



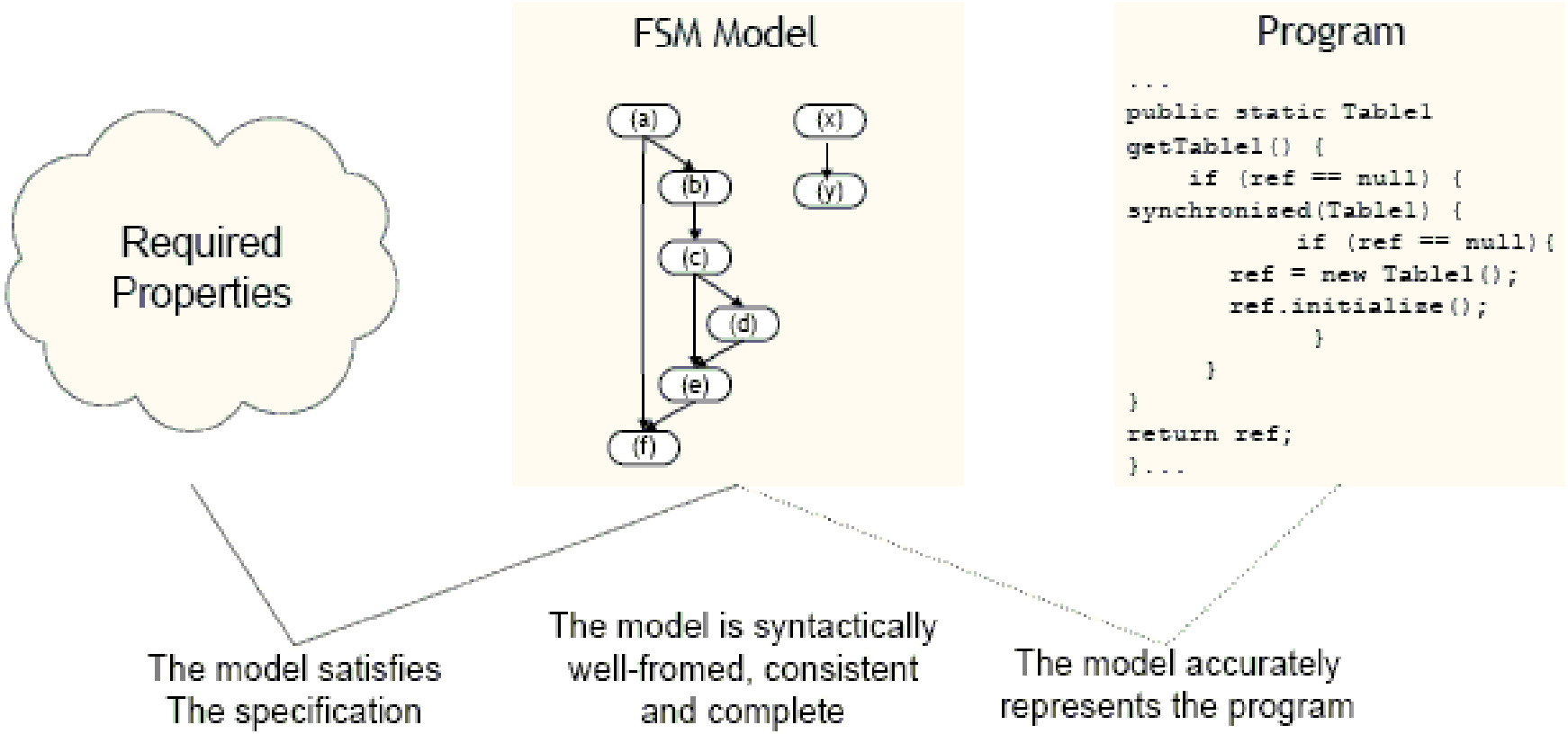
Tabular representation

	LF	CR	EOF	other
e	e / emit	l / emit	d / –	w / append
w	e / emit	l / emit	d / emit	w / append
l	e / –		d / –	w / append

LF: line feed
CR: carriage return
EOF: end-of-fine
other: everything else
(e.g., d / emit means "emit and then proceed to state d")

Note. In this FSM a transition from state *l* on a CR event is omitted.

Using Models to Reason about System Properties

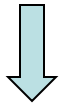


```
1  /** Convert each line from standard input */
2  void transduce() {
3
4     #define BUFLEN 1000
5     char buf[BUFLEN];   /* Accumulate line into this buffer   */
6     int   pos = 0;          /* Index for next character in buffer */
7
8     char inChar; /* Next character from input */
9
10    int atCR = 0; /* 0="within line", 1="optional DOS LF" */
11
12    while ((inChar = getchar()) != EOF ) {
13       switch (inChar) {
14       case LF:
15          if (atCR) {    /* Optional DOS LF */
16             atCR = 0;
17          } else {        /* Encountered CR within line */
18             emit(buf, pos);
19             pos = 0;
20          }
21          break;
22       case CR:
23          emit(buf, pos);
24          pos = 0;
25          atCR = 1;
26          break;
27       default:
28          if (pos >= BUFLEN-2) fail("Buffer overflow");
29          buf[pos++] = inChar;
30       } /* switch */
31    }
32    if (pos > 0) {
33       emit(buf, pos);
34    }
35 }
```

Model Abstraction

is used building an FSM from a program

Abstract state	*Concrete state*		
	Lines	atCR	pos
e (Empty buffer)	2 – 12	0	0
w (Within line)	12	0	> 0
l (Looking for LF)	12	1	0
d (Done)	35	–	–



	LF	CR	EOF	other
e	e / emit	l / emit	d / –	w / append
w	e / emit	l / emit	d / emit	w / append
l	e / –	l / emit	d / –	w / append

Note. An FSM is generated immediately from a tabular representation.

Summary

- Models must be much simpler than the artifact they describe to be understandable and analyzable
- Must also be sufficiently detailed to be useful
- CFG are built from software
- FSM can be built before software to documentintended behavior